

# 1. History & Introduction

- Dennis Ritchie C99.

- 注释  
// 分行 /\*

- infix notation 1+3.

- 没有 exponentiation operator (e.g.  $x^n$ )

- rounds towards 0 (向下取整)

- trace -X (exp) 路径追踪

int. long. bool. char. double string. ptr symbol

- entry point : main (return 0)

- Bool	F : 0	T : 1	
	Character	char	%c
	float	float	%f
	double	(比 float 更精确)	

a. 97  
("%c", a) ≠ ("%d", a)

- Control flow {  
1. compound ( {...} )  
2. function call  
3. conditionals (if)  
4. iteration (loop)

\* return statement ends the function

- "=" & "=="

"=" 赋值 initialization

"!=" / "==" 判断是否相等 (equality operator)

$x=y \neq y=x$

static type system.

identifier (function name)

top level expression

```
int my-add (int a, int b) {
    return a+b;
}
my-add { 1, 2 } arguments
```

```
8 int n = 5; // initialization syntax
9 n = 6; // assignment operator
```

Example 1.5.5: dangling else

```
1 // this program demonstrates the dangers of a "dangling else"
2
3 #include "cs136.h"
4
5 // win_or_lose(y) ... what DOES this function do ??
6 void win_or_lose(int y) {
7     printf("win_or_lose(%d):\n", y);
8     if (y > 0)
9         if (y != 7) {
10            printf("you lose\n");
11        }
12     else {
13         printf("you win!\n"); } // when does this print?
14 }
15 int main(void) {
16     win_or_lose(-1); // no output
17     win_or_lose(1);
18     win_or_lose(7);
19 }
```

Code Output

```
win_or_lose(-1):
win_or_lose(1):
you lose
win_or_lose(7):
you win!
main.c:11:3: warning: add explicit braces to avoid dangling else [-Wdangling-else]
    else
    ^
1 warning generated.
```

← 不输出. 因为默认 {}

## 2. Imperative C

- functional programming paradigm
  - ① "mathematical" = "pure"
    - only return values
    - answer depend on argument values
  - ② only constants
- imperative programming paradigm
  - ① function 不 return values
  - ② 使用 variables & constants
  - ③ {...} compound statement
  - ④ side effects

- 3 types of statement
    - compound ~ {...} 接收序执行
    - expression ~ 目的: generate side effects (值被丢掉)
    - control flow ~ (return, if, else) 控制执行顺序
- not all expression statement generate side effect

### - side effects

If num passed to function. It output a message and returns num.

若 function 除了输出 value 时, 还产生 output. 则 function 有 side effect.

- \* assert & trace-int 不产生 side effects
- \* function may only have side-effects
- \* 所有有 side-effect in function 为 impure function

① 0 (produce output)  
printf(" ").

```
printf("2 plus 2 is: %d\n", 2 + 2);
```

format specifier. →

整数	%d
换行	\n
斜杠	\/
双斜杠	\/\
双引号	\"

```
printf("five\n")
```

return value 为 5. (字符数)  
output 为 "five"

\* %04d 一共四个数字. 其余用 0 补齐  
eg. 0042.

② I (read input)

int n = read\_int();

③ Mutation (change the value of variables)

- mutable global variable

改变 variable in 它 ep. 不断定义 m = ...

(constant is immutable)

A global mutable variable is "impure" even if it has no side effects  
- ep n+=1

- mutate a variable through pointer parameter

~~✗~~ mutating a local variable/parameter don't have side effects  
Since it does not affect state outside of the function

immutable → constant. (const)

mutable → variable { global scope 为全部  
local "block scope": 从 initialize 到 函数结尾  
↑  
(identifier 前不带 const 都是 variable)

### Testing Terminology

Now that we've learned more about testing, we can summarize testing terminology:

**Assertion testing:** using assertions to test our code    test return value  
**I/O-driven testing:** uses input and expected output    test output  
**Testing harness:** using input to call functions

**White box testing:** you can see the code being tested

**Black box testing:** you can not see the code being tested

**Unit testing:** testing one piece at a time

Sometimes the changes fix the original bug but introduce other bugs, breaking code that previously worked. In this case the code is said to have "regressed" (gotten worse).

**Regression testing** reruns all tests to check if everything still works (after a change to any part of the code) to ensure that changes made don't introduce new bugs.

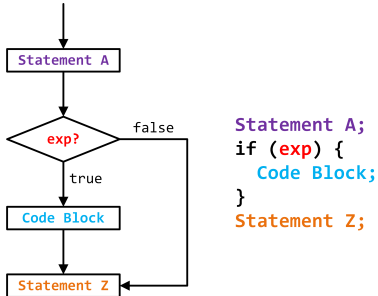
# 3. Memory and Control Flow

- Control flow. program 如何执行.

- types {
- compound statement (block)
  - function call
  - conditional (if / else)
  - iteration (loops)

需要 mutation (side-effect) 使循环停止

if

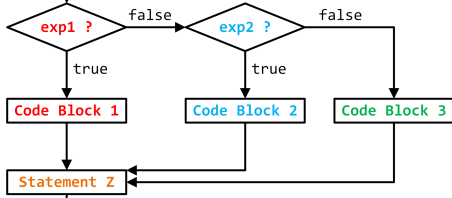


```

Statement A;
if (exp) {
  Code Block;
}
Statement Z;
  
```

exp 相互无关

else if

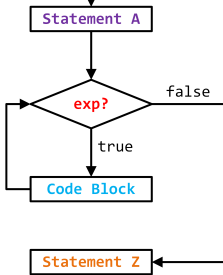


```

Statement A;
if (exp1) {
  Code Block 1;
} else if (exp2) {
  Code Block 2;
} else {
  Code Block 3;
}
Statement Z;
  
```

exp 相互有关

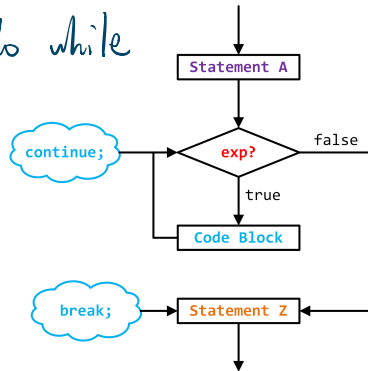
while



```

Statement A;
while (exp) {
  Code Block;
}
Statement Z;
  
```

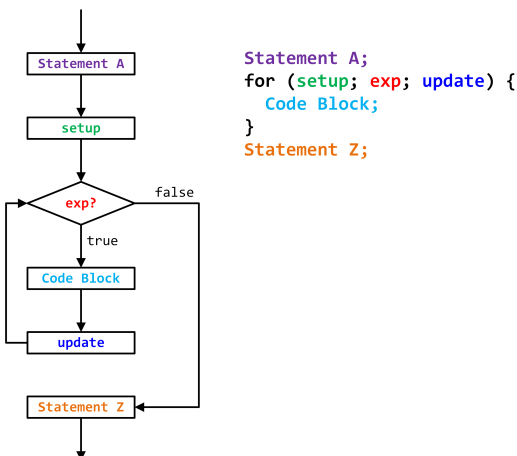
do while



```

Statement A;
while (exp) {
  Code Block;
}
Statement Z;
  
```

for



```

Statement A;
for (setup; exp; update) {
  Code Block;
}
Statement Z;
  
```



simple recursion

```
int recursive_sum(int k) {
    if (k <= 0) {
        return 0;
    }
    return k + recursive_sum(k - 1);
}
```

while loop

```
int iterative_sum(int k) {
    int s = 0;
    while (k > 0) {
        s += k;
        --k;
    }
    return s;
}
```

accumulative recursion

```
int accsum(int k, int acc) {
    if (k <= 0) {
        return acc;
    }
    return accsum(k - 1, k + acc);
}
int recursive_sum(int k){
    return accsum(k, 0);
}
```

while loop

```
int iterative_sum(int k) {
    int acc = 0;
    while (k > 0) {
        acc += k;
        --k;
    }
    return acc;
}
```

do while

while

```
int main(void) {
    int i = 0;

    printf("Do..while loop:\n");
    // do...while
    i = 5; // setup
    do {
        printf("%d\n", i); // statement
        --i; // update
    } while (i >= 0); // expression

    printf("while loop:\n");
    // while
    i = 5; // setup
    while (i >= 0) { // expression
        printf("%d\n", i); // statement
        --i; // update
    }
}
```

在刚开始就不满足 while 后条件的情况下，do...while will do 1 more iteration than while

while loop

for loop

```
int main(void) {
    int i = 0;

    printf("While loop:\n");
    i = 5; // setup
    while (i >= 0) { // expression
        printf("%d\n", i);
        --i; // update
    }

    printf("For loop:\n");
    for (i = 5; i >= 0; --i) {
        printf("%d\n", i);
    }
}
```

返回

5

4

3

2

1

0

set up statement.

```
while (expression) {
    body;
    update (循环)
}
```

```
for (setup; expression; update) {
    body;
}
```

recursion 在 function 中不断调用自己

iteration 用 loops.

efficiency 相同

## - Memory.

one bit of storage has 2 states: 0 1.

byte: smallest accessible unit of memory.

address: position of memory.

$$2^{\text{bytes}} = \text{bits}$$

identifier	type	# bytes	starting address
n	int	4	0x5000

variable definition 的作用: ① find space to store variable  
② track address  
③ 将值存进 address

size operator: (sizeof) 是 operator.

int: 4 bytes (32 bits)  $2^{32}$  possible values can represent

If you have a machine that uses 32-bit addresses, what is the maximum amount of RAM that can be addressed?

4GB (4,294,967,296 bytes, i.e.  $2^{32}$ )

$$2^{32} \rightarrow 2^{16} \rightarrow 2^8 \rightarrow 2^4$$

overflow 溢出: 可生成任何值

## - Sections of memory.

Code
Read-Only Data
Global Data
Heap
Stack

↔ the call stack is stored

# - Structure struct

```
4
5 struct posn {           // name of the structure
6     int x;             // type and field names
7     int y;
8 };                    // don't forget this ;
9
10 // posn_equal(a, b) returns true the structs a and b are equal and
11 // false otherwise
12 bool posn_equal (struct posn a, struct posn b) {
13     return (a.x == b.x) && (a.y == b.y);
14 }
15
16 int main(void) {
17     struct posn p = {1, 2};
18
19     p.x = 5;           // VALID MUTATION
20     p.y = 6;           // VALID MUTATION
21
22     // alternatively:
23     struct posn new_p = {5, 6};
24     p = new_p;         // VALID MUTATION
25
26     trace_int(p.x);
27     trace_int(p.y);
28
29     if (posn_equal (p, new_p)) {
30         printf("The structures p and new_p are equal!\n");
31     }
32
33     printf("The value of p is (%d, %d)\n", p.x, p.y);
34 }
```

== & !=

不应用于 structure

## Code Output

```
The structures p and new_p are equal!
The value of p is (5, 6)
>>> [main.c|main|26] >> p.x => 5
>>> [main.c|main|27] >> p.y => 6
```

class exercise 3-1 ~~4~~

# - Stacks & recursion



从下往上存

stack frame {  
 argument value  
 local variable  
 return address

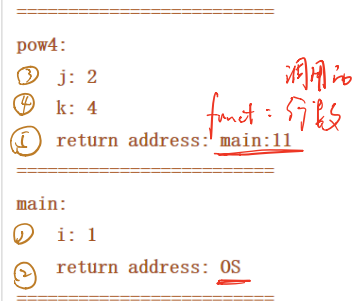
before function can be called.  
 all arguments must be values

return address : 调用此函数前程序的地址.

main 的 return address : OS (operating system)

```

1 int pow4(int j) {
2     printf("inside pow4\n");
3     int k = j * j;
4     // Snapshot
5     return k * k;
6 }
7
8 int main(void) {
9     printf("inside main\n");
10    int i = 1;
11    printf("%d\n", pow4(i + i));
12 }
    
```

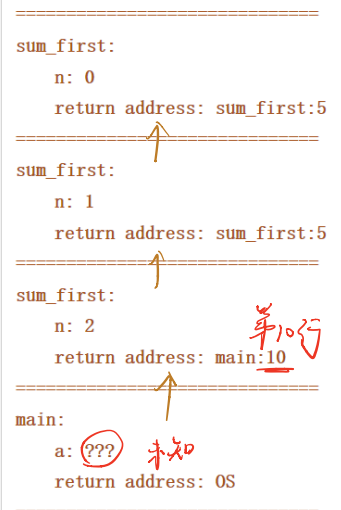


global variable 在执行前保留  
 local variable 在调用函数时保留

## - recursion

```

1 int sum_first(int n) {
2     if(n == 0) { // Snapshot here
3         return 0;
4     } else {
5         return n + sum_first(n - 1);
6     }
7 }
8
9 int main(void) {
10    int a = sum_first(2);
11    //...
12 }
    
```



# 4. Pointer

- pointer: variable that stores a memory address %p

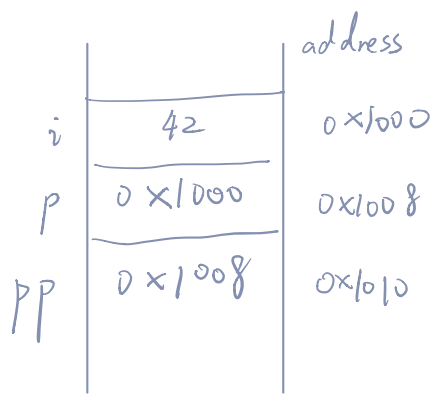
\* : 定义 pointer

& : 赋值 address operator produce location of an identifier

```
int *p = &i
```

pointer ∈ variable → store value (address)

size of pointer. 64 bits (8 bytes)

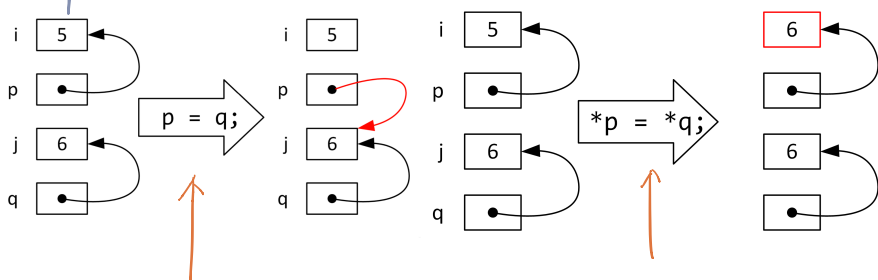


```
printf ("%p", &p)
```

```
printf ("%d", *p)
```

```
int **pp = &p.
```

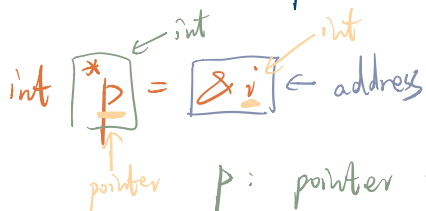
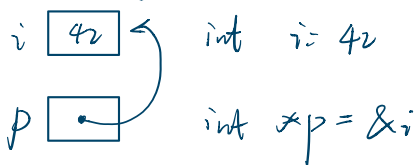
pointer to pointer



p 指向 q 指向 i  
p 的值改变 (p=q)

p 指向 q 指向的值  
p 指向的值改变。  
(value = q 指向的)

## - memory diagram



p: pointer to an integer i = i in address

\*p: p 里所存的值 = i 的值

p 与 &i return 相同

indirection operator

- 在 Structure 中出现 pointer.

(.) 优先级高于 (\*)

ptr → field ≡ (\*) ptr . field

↑  
indirection selection operator

```
5 struct posn {
6     int x;
7     int y;
8 };
9
10 int main(void) {
11     struct posn my_posn = {0, 0};
12     struct posn *ptr = &my_posn;
13
14     *ptr.x = 3; // awkward
15     ptr->y = 4; // 相当于 * + . // much better
16
17     trace_int(ptr->x);
18     trace_int(ptr->y);
19 }
```

- Pointer assignment 指针赋值  
aliasing: 多个指针访问同一地址

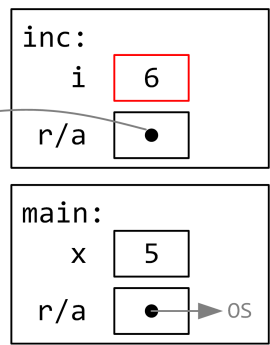
```
i = 1;
int *p1 = &i;
int *p2 = p1;
int **p3 = &p1;
```

\*p3, p2, p1  
都访问 i 的地址

- Mutation

```
void inc(int i) {
    ++i;
}

int main(void) {
    int x = 5;
    inc(x);
}
```

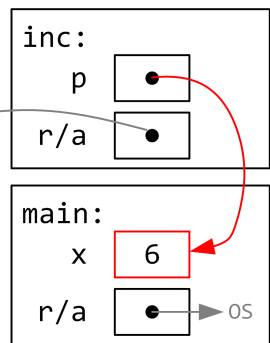


使用 pointer 不会破坏原始值  
x 永远 = 5.

```
void inc(int *p) {
    *p += 1;
}

int main(void) {
    int x = 5;
    inc(&x);
}
```

x=5  
x=6  
x=7  
...



inc(&x) 相当于地址所对应的值

side effect → mutate a variable through a pointer parameter

Example 4.2.3: mutation side effects

```
1 // This program illustrates a "swap" function in C
2 // that uses pointers
3
4 #include "cs136.h"
5
6 // swap(px, py) "swaps" the contents of *px and *py
7 // effects: modifies *px and *py
8 // requires: px and py are valid pointers
9 void swap(int *px, int *py) {
10     assert(px);
11     assert(py);
12     int temp = *px;
13     *px = *py;
14     *py = temp;
15 }
16
17
18 int main(void) {
19     int a = 3;
20     int b = 4;
21     trace_int(a);
22     trace_int(b);
23     swap(&a, &b); // Note the &
24     trace_int(a);
25     trace_int(b);
26 }
```

检测 \*px 与 \*py 中存的值 不是 NULL  
↑  
空值 & value

```
Code Output
>>> [main.c|main|21] >> a => 3
>>> [main.c|main|22] >> b => 4
>>> [main.c|main|24] >> a => 4
>>> [main.c|main|25] >> b => 3
```

function 中不能 return pointer. (如果一开始没有定义的话)

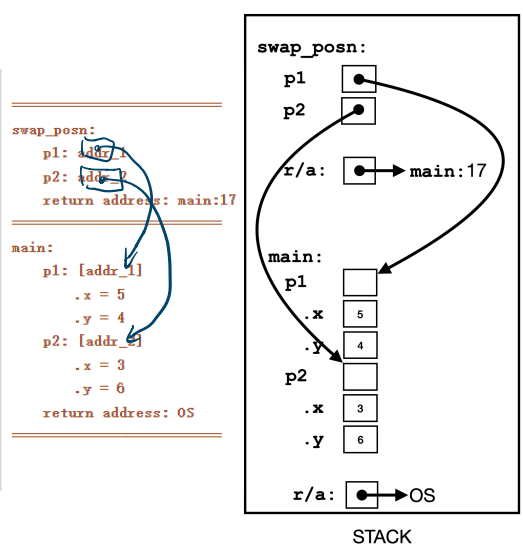
```
1 // This program demonstrates returning a pointer (syntax) and
2 // returning a pointer to a stack frame
3
4 #include "cs136.h"
5
6 // bad_idea(n) returns a pointer to the bad_idea frame (BAD!)
7 int *bad_idea(int n) {
8     return &n; // NEVER do this
9 }
10
11 // ptr_to_max(a, b) returns either a or b: whichever points to
12 // the larger value
13 // requires: a, b are valid pointers
14 int *ptr_to_max(int* a, int* b) {
15     assert(a);
16     assert(b);
17     if (*a >= *b) {
18         return a;
19     }
20     return b;
21 }
22
23 int main(void) {
24     int x = 3;
25     int y = 4;
26
27     int *p = ptr_to_max(&x, &y); // note the &
28     assert(p == &y);
29     printf("%p\n", bad_idea(13));
30 }
```

Code Output

```
0x7f05ca100060
main.c:8:11: warning: address of stack memory associated with parameter 'n' returned [-Wreturn-stack-address]
    return &n;
           ^
1 warning generated.
```

stack frames with pointers

```
1 void swap_int(int *i, int *j) {
2     int temp = *i;
3     *i = *j;
4     *j = temp;
5 }
6
7 void swap_posn(struct posn *p1,
8               struct posn *p2) {
9     swap_int(&(p1->x), &(p2->x));
10    // Snapshot
11    swap_int(&(p1->y), &(p2->y));
12 }
13
14 int main(void) {
15     struct posn p1 = {3, 4};
16     struct posn p2 = {5, 6};
17     swap_posn(&p1, &p2);
18 }
```





## - Input & Pointer

`scanf("%d", &i)` → return value : int.

成功读到值 或 EOF (end of file)

side effects of scanf:

- ① a value is read from input
- ② i is mutable
- ③ ... is mutated

### ✱ 4.3.1

```
5 // read_sum() reads ints from input (until failure)
6 // and returns their sum
7 // effects: reads input
8 int read_sum(void) {
9     int sum = 0;
10    int n = 0;
11    while (scanf("%d", &n) == 1) {
12        sum += n;
13    }
14    return sum;
15 }
16
17 int main(void) {
18     printf("%d\n", read_sum());
19 }
```

用于判断是否还有值

### ✱ 4.3.2

```
14 int readint(void) {
15     int i = 0;
16     int result = scanf("%d", &i);
17     if (result == 1) {
18         return i;
19     }
20     return READINT_FAIL;
21 }
```

读取值 in type

将读取的值存入 &i 中。

```
int count_char(void) {
    int count = 0;
    char c = '\0';
    while (scanf("%c", &c) == 1) {
        count++;
    }
    return count;
}
```

"%c" 不忽略空格

"%c" 忽略空格

### ✱ 4.3.5

## - Structure & pointer

原因: ① avoid copying the structure  
② mutate the contents of the structure

<code>int *p;</code>	p can point at any mutable integer, you can modify the int (via *p)
<code>const int *p;</code>	p can point at any integer, you can NOT modify the integer (via *p)
<code>int * const p = &amp;i;</code>	p always points at the integer i, i must be mutable and can be modified (via *p)
<code>const int * const p = &amp;i;</code>	p always points at the integer i, you can not modify i (via *p)

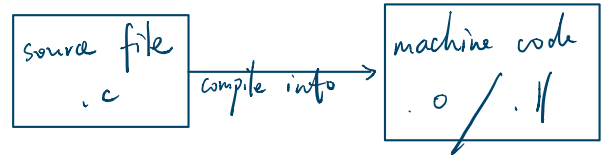
## - function pointers

```
5 // io_apply(f) reads in each int [n] from input
6 // and prints out f(n)
7 // requires: f is a valid pointer
8 // effects: produces output
9 // reads input
10 void io_apply(int (*f)(int)) {
11     assert(f);
12     int n = 0;
13     while (scanf("%d", &n) == 1) {
14         printf("%d\n", f(n));
15     }
16 }
17
18 // sqr(i) calculates i^2
19 int sqr(int i) {
20     return i * i;
21 }
22
23 int main(void) {
24     io_apply(sqr);
25 }
```

return-type (\* function-name) (parameter-type ...)

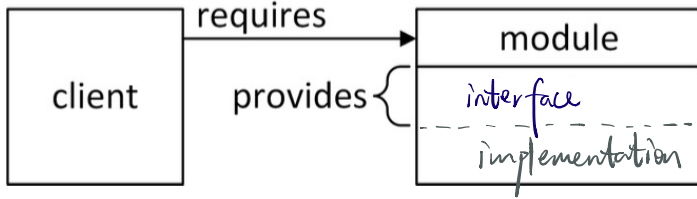
# 5. Modularization

- def. modularization



divide programs into well defined modules

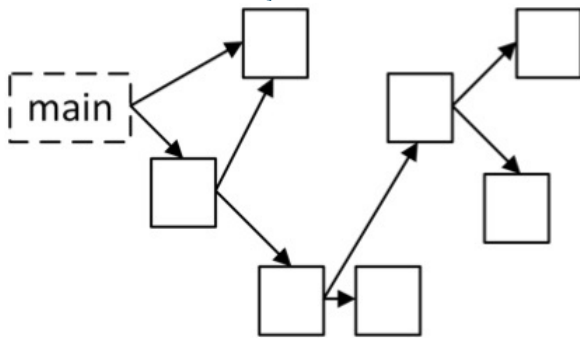
→ "client" (main.c) requires the functions that a "module" (fun.c) provides.



≡ libraries

functions in same module should share same purpose

→ module dependency graph. \* no cycles. \* have "roots" (main)



- properties
- reusability 可同时用作...
  - abstraction 省略. 简洁
  - readability 文档齐全
  - maintainability 可维护

## Concept Check 5.1.4

2.0/2.0 points (graded)

Suppose we have developed a module containing functions that model sunlight reflecting off of moving water (like a river or stream). For each term choose the number for the statement below that best illustrates an example of this term.

1. A client can use the module without understanding the physics used in the calculations

2. The module is well documented readability

3. The module is used in both video games and animated movies

4. Some functions in the module can be updated to execute more efficiently without changing the interface

Re-usability

3 ✓

Abstraction

1 ✓

outside

Readability

2 ✓

Maintainability

4 ✓

Show answer

- definition includes declaration

int foo (int x); → declaration

int foo (int x) {  
...  
} → definition

```
// Declaration (odd.h):  
bool is_odd(int number);
```

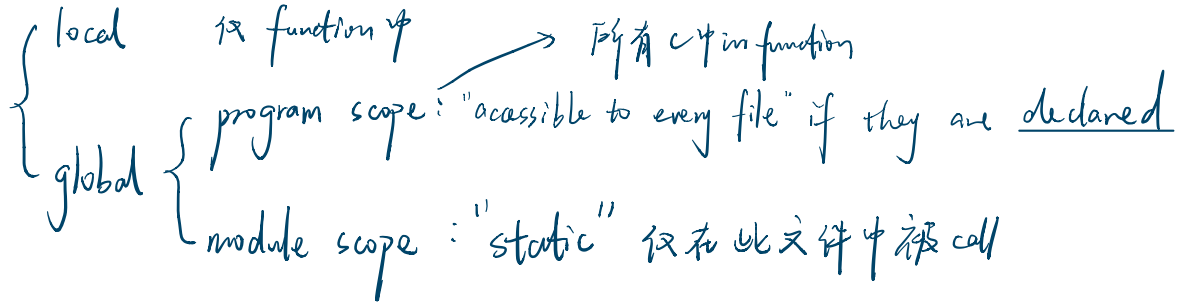
```
// Definition (odd.c):  
bool is_odd(int number) {  
    return (number % 2 != 0);  
}
```

variable  
declaration  
可跨文件

```
1 // this program demonstrates VARIABLE declarations  
2  
3 #include <stdio.h>  
4 #include "cs136-trace.h"  
5  
6 extern int g; // variable DECLARATION  
7  
8 int main(void) {  
9     printf("%d\n", g); // this is now ok  
10    trace_int(g);  
11 }  
12  
13 int g = 7; // variable DEFINITION
```

```
main.c  module.c  + Add File  
1 // main.c  
2 // This program demonstrates PROGRAM scope  
3 // with a variable declaration  
4  
5 #include <stdio.h> ←  
6 #include "cs136-trace.h" ← " "  
7  
8 // NEW declaration 在 - } directory 中  
9 extern int g;  
10  
11 // g IN scope  
12  
13 int main(void) {  
14     printf("%d\n", g); // ok  
15 }
```

- type of scope



- module interface (placed in ... .h) header files

def. list of function that module provides

- implementation (.. .c) 分析
- provided to clients

- hidden from client
- information hiding.
  - hide implementation details (provide machine code .. .ll instead .. .c)
  - hiding data. { security flexibility

- 包含 {
1. overall description
  2. function declaration.
  3. documentation (esp. purpose)

... .c won't compile if u don't include ... h

C isn't good at hiding data

若 b.c 中 include a.h.  
则 b.h 也要写 #include "a.h"

## - #include

a preprocessor directive temporarily "modifies" file before it is run  
在run之前提前预处理

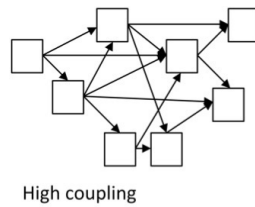
< > specify module which is one of the standard modules (C provide)  
" " "regular modules" (自定义的) (ep. libraries)

<assert.h>	provides the function assert
<limits.h>	provides the constants INT_MAX and INT_MIN
<stdbool.h>	provides the bool data type and the constants true and false
<stdlib.h>	provides the constant NULL and abs()

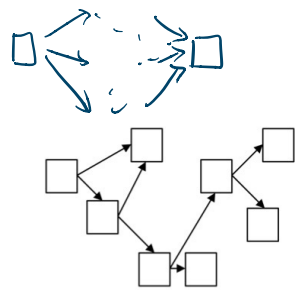
## - cohesion & coupling

high cohesion : all interface functions are related & working towards a common goal

low coupling : little interaction between modules



High coupling



Low coupling

opaque structure : module only provides incomplete declaration in interface

↑  
incomplete declaration : only pointers to structure can be defined  
\* supply a complete definition of structure in .c file

transparent structure : client can access and use without pointer.  
put complete definition of structure in .h file

- ADT (abstract data type) 数据结构

def. mathematical model for storing and accessing data through operations.  
implemented as data storage modules.  
implemented with data structure

Collection ADT: an ADT designed to store an arbitrary number of items.

- 数据结构 data structure

- dictionary ADT (pairs of keys & values)

lookup  
insert  
remove

- stack ADT

push  
pop  
top  
is-empty

- queue ADT

add-back  
remove-front  
front  
is-empty

- sequence ADT

item-at  
insert-at  
remove-at  
length.

- set ADT

add  
remove  
length  
member

⊗ class exercise

# b. Arrays

b.2.1 ?

- Compound data storage : structure & array.
- def. can create array of any built-in data type in C  
a data structure that contain fixed number of element all same type

## - array initialization

uninitialized global array are zero filled.  
uninitialized local array are filled with arbitrary values from stack ~~ptr~~

- length  $a[n] = \{ \dots \}$

size bytes num =  $\frac{\# \text{ element}}{n} \times \# \text{ size}$

```
1 // This program demonstrates how the value of an array
2 // is the same as the address of the first element
3
4 #include "cs136-trace.h"
5
6 int main(void) {
7     int a[6] = {4, 8, 15, 16, 23, 42};
8
9     trace_ptr(a);
10    trace_ptr(&a);
11    trace_ptr(&a[0]);
12    trace_ptr(&a[1]);
13
14    trace_int(a[0]);
15    trace_int(*a);
16 }
```

array doesn't have values

## Code Output

```
>>> [main.c|main|9] >> a => 0x7f691b200020
>>> [main.c|main|10] >> &a => 0x7f691b200020
>>> [main.c|main|11] >> &a[0] => 0x7f691b200020
>>> [main.c|main|12] >> &a[1] => 0x7f691b200024
>>> [main.c|main|14] >> a[0] => 4
>>> [main.c|main|15] >> *a => 4
```

# - array & pointers

\* pointer arithmetic :  $p \pm i = p \pm i \times \text{size of } (*p)$

若两个  $p$  point to same type, 则可以相减

$$p_1 - p_2 = \frac{p_1 - p_2}{\text{size of } (*p)}$$

$$a[0] \equiv *a$$

\* b.2.1

\* b.2.2

\* b.3.3

```
8
9 int sum_array(const int *a, const int len) {
10     assert(a);
11     assert(len > 0);
12     int sum = 0;
13     for (const int *p = a; p < a + len; ++p) {
14         sum += *p;
15     }
16     return sum;
17 }
18
19 // traditional array notation
20
21 int sum_array_old(const int a[], const int len) {
22     assert(a);
23     assert(len > 0);
24     int sum = 0;
25     for (int i = 0; i < len; ++i) {
26         sum += a[i];
27     }
28     return sum;
29 }
```

← 不是 a. 因为 a 不会改变



- quick sort 排序

```

6
7 // quick_sort_range(a, first, last) sorts the elements of
8 // a in the range a[first]..a[last] (inclusive)
9 // in ascending order
10 // requires: a is a valid array in the range first..last [not as
11 // effects: modifies a
12 void quick_sort_range(int a[], int first, int last) {
13     assert(a);
14     if (last <= first) return; // length is <= 1
15
16     int pivot = a[first]; // first element is the pivot
17     int pos = last; // where to put next larger
18
19     for (int i = last; i > first; --i) {
20         if (a[i] > pivot) {
21             swap(&a[pos], &a[i]);
22             --pos;
23         }
24     }
25     swap(&a[first], &a[pos]); // put pivot in correct place
26     quick_sort_range(a, first, pos - 1);
27     quick_sort_range(a, pos + 1, last);
28 }
29
30 // quick_sort(a, len) sorts the elements of a in ascending order
31 // requires: len > 0
32 // effects: modifies a
33 void quick_sort(int a[], int len) {
34     assert(a);
35     assert(len > 0);
36     quick_sort_range(a, 0, len - 1);
37 }
38
39 int main(void) {
40     int a[7] = {8, 6, 7, 5, 3, 0, 9};
41     print_array(a, 7);
42     quick_sort(a, 7);
43     print_array(a, 7);
44 }

```

Code Output

```

8, 6, 7, 5, 3, 0, 9.
0, 3, 5, 6, 7, 8, 9.

```

- binary search 找数

Example 6.3.5: Binary Search

```

main.c - Add File
1 // This program demonstrates a binary search
2 // on a sorted array
3
4 #include <assert.h>
5 #include "cs136-trace.h"
6
7 // find_sorted(item, a, len) finds the index of item in a,
8 // or returns -1 if it does not exist
9 // requires: a is sorted in ascending order [not asserted]
10 // len > 0
11 int find_sorted(int item, const int a[], int len) {
12     assert(a);
13     assert(len > 0);
14     int low = 0;
15     int high = len - 1;
16     while (low <= high) {
17         int mid = low + (high - low) / 2;
18         if (a[mid] == item) {
19             return mid;
20         } else if (a[mid] < item) {
21             low = mid + 1;
22         } else {
23             high = mid - 1;
24         }
25     }
26     return -1;
27 }
28
29 int main(void) {
30     int a[10] = {1, 2, 3, 5, 7, 11, 13, 17, 19, 23};
31
32     assert(find_sorted(1, a, 10) == 0);
33     assert(find_sorted(13, a, 10) == 6);
34     assert(find_sorted(23, a, 10) == 9);
35     assert(find_sorted(8, a, 10) == -1);
36 }

```

↑  
↑ 适用于 sorted list  
↑ use "divide and conquer"

- Oversized Array (len 过大)

```
int arr[10] = {115, 116, 136, 245, 246};
```

Select the correct dropdown option.

```
printf("%d", arr[6]);
```

0

```
printf("%d", arr[0]);
```

115

```
printf("%d", arr[-1]);
```

An error/warning or unexpected behaviour will occur due to invalid array index.

```
printf("%d", arr[136]);
```

An error/warning or unexpected behaviour will occur due to invalid array index.

## 7. Efficiency

- algorithm. step-by-step description of how to solve a problem.

- efficiency.

time  $\sim$ : how long it takes to solve a problem.

space  $\sim$ : how much memory requires to solve a problem.

power  $\sim$

- running time  $T(n)$

measure the number of elementary operators required.

$n$ : size/length of data

worst case running time

- Big O

"order" of running time: "dominant" term. without constant coefficient.

2 algorithm are equivalent if they have same order.

growth rate (快)

↓

$O(1)$

$O(\log n)$

$O(n)$

$O(n \log n)$

$O(n^2)$

$O(n^3)$

↘  $O(2^n)$

加法: result = 较大数

$$O(\log n) + O(n) = O(n)$$

乘法: result = 两者相乘

$$O(\log n) \times O(n) = O(n \log n)$$

# - 计算

## ① Simple

```
// max(a, b) returns the max of the elements
// time: O(1)+O(1)+ ... +O(1) = O(1)
int max(int a, int b) {
    if (a > b) {
        return a;
    } else {
        return b;
    }
}
```

## ② Iterative analysis (summations)

1. 内层  $\rightarrow$  外层
2. worst case 迭代次数. 数据大小
3.  $T(n)$

$$\sum_{i=1}^{n/10} O(1) = \sum_{i=1}^n O(1)$$

ep. 

```
for (i = 1; i <= n; ++i) {
    printf("*");
}
```

$$T(n) = \sum_{i=1}^n O(1) = O(1) \times n = O(n)$$

$$\begin{aligned} \sum_{i=1}^{\log n} O(1) &= O(\log n) \\ \sum_{i=1}^n O(1) &= O(n) \\ \sum_{i=1}^n O(n) &= O(n^2) \\ \sum_{i=1}^n O(i) &= O(n^2) \\ \sum_{i=1}^n O(i^2) &= O(n^3) \end{aligned}$$

ep. 

```
for (i = 0; i < n; ++i) {
    for (j = 0; j < i; ++j) {
        printf("*");
    }
    printf("\n");
}
```

$\leftarrow$  nested loops

$$\text{inner loop: } \sum_{j=0}^{i-1} O(1) = O(i)$$

$$\text{outer loop: } \sum_{i=0}^{n-1} (O(i) + O(i)) = O(n) + O(n^2) = O(n^2)$$

外层      内层

## ③ Recursion

1. order of function. (不包括 recursion)
2.  $T$  递归 in data size
3. recurrence relation
4. close-form sol.

```
int sum_first(int n) {
    if (n == 0) {
        return 0;
    } else {
        return n + sum_first(n - 1);
    }
}
```

1. non-recursive operators:  $O(1)$ ,  $+$ ,  $-$ ,  $=$ .
2.  $n-1$
3.  $T(n) = O(1) + T(n-1)$
4.  $T(n) = O(n)$

The recurrence relations we encounter in this course are:

$T(n) = O(1) + T(n - k_1)$	$= O(n)$
$T(n) = O(n) + T(n - k_1)$	$= O(n^2)$
$T(n) = O(n^2) + T(n - k_1)$	$= O(n^3)$
$T(n) = O(1) + T(\frac{n}{k_2})$	$= O(\log n)$
$T(n) = O(1) + k_2 * T(\frac{n}{k_2})$	$= O(n)$
$T(n) = O(n) + k_2 * T(\frac{n}{k_2})$	$= O(n \log n)$
$T(n) = O(1) + T(n - k_1) + T(n - k_1')$	$= O(2^n)$

where  $k_1, k_1' \geq 1$  and  $k_2 > 1$

## ★ - Sorting

(same efficiency)

Algorithm	best case	worst case
selection sort	$O(n^2)$	$O(n^2)$
insertion sort	$O(n)$	$O(n^2) > \sum_{i=1}^n \sum_{j=1}^i O(1) = O(n^2)$
quick sort	$O(n \log n)$	$O(n^2)$

## - Binary Search

$$T(n) = \sum_{i=1}^{\log_2 n} O(1) = O(\log n)$$

$n$  为 loop 次数

```
void f4(int n) {
    for (int i = 0; i < n; i += 2) {
        for (int j = n; j > 0; j /= 2) { }
        printf("*");
    }
    printf("\n");
}
```

inner loop:  $\sum_{i=1}^{\log n} O(1) = O(\log n)$

outer loop:  $\sum_{i=1}^n O(\log n) + O(1)$   
 $= O(n \log n) + O(n)$   
 $= O(n \log n)$

## 8. String

- C string: array of char. terminated by null character.

↑  
terminator '\0' = 0

- initialization initial in 位置输入 '\0'

- null terminated

while (s[i]) 检测 array 是否结束

- string len int strlen(s)

```
// my_strlen(s) behaves the same as the built-in strlen function
// the length does not include the null character
// time: O(n)
int my_strlen(const char s[]) {
    assert(s);
    int len = 0;
    while (s[len]) {
        ++len;
    }
    return len;
}

// my_strlen_ptr(s) behaves the same as the built-in strlen
// the length does not include the null character
// time: O(n)
int my_strlen_ptr(const char *s) {
    assert(s);
    const char *p = s;
    while (*p) {
        ++p;
    }
    return (p - s);
}
```

- lexicographical order int strcmp(s1, s2)

```
// my_strcmp(s1, s2) behaves the same as strcmp and returns:
// s1 precedes s2 => negative int
// s1 identical to s2 => 0
// s1 follows s2 => positive int
// time: O(n), n is min of the lengths of s1, s2
int my_strcmp(const char s1[], const char s2[]) {
    assert(s1);
    assert(s2);
    int i = 0;
    while (s1[i] == s2[i] && s1[i]) {
        ++i;
    }
    return s1[i] - s2[i];
}

int main(void) {
    char a[] = "the same?";
    char b[] = "the same?";
    char c[] = "the samejQuery22408322300989689644_1654783296792";
    char *s = a;

    assert(a != b);
    assert(a == s);

    assert(my_strcmp(a, b) == 0);
    assert(my_strcmp(a, c) < 0); // -43
    assert(my_strcmp(c, a) > 0); // 43
}
```

("ab", "a") → 1

("a", "b") → -1

("", "ab") → -1

- strcpy char \*strcpy (char \*dest, const char \*src)  
将 dest 与 src 换位

```
// my_strcpy(dest, src) behaves the same as the built-in strcpy function
// time: O(n), n is the length of src
char *my_strcpy(char *dest, const char *src) {
    assert(dest);
    assert(src);
    char *d = dest;
    while (*src) {
        *d = *src;
        ++d;
        ++src;
    }
    *d = '\0';
    return dest;
}

int main(void) {
    char a[] = "cat";
    char b[4] = "dog";
    my_strcpy(b, a);
    trace_string(a);
    trace_string(b);
    assert(!strcmp(a, b)); // !a==b
}
```

buffer overflow (overrun) ∴ 确保 dest array 足够大  
↳ arr[2] = "hi" 数组大小为3. 'h', 'i', '\0'  
致溢

- strcat char \*strcat (char\* dest, const char\* src)  
将 src 加入到 dest 后面 (合并两个 char)

```
// my_strcat(dest, src) behaves the same as the built-in strcat function
// time: O(n + m) n,m are lengths of src,dest
char *my_strcat(char *dest, const char *src) {
    assert(dest);
    assert(src);
    strcpy(dest + strlen(dest), src);
    return dest;
}

int main(void) {
    char a[] = "cat";
    char b[7] = "dog";
    my_strcat(b, a);
    trace_string(a);
    trace_string(b);
    assert(!strcmp(b, "dogcat"));
}
```

array comparison always evaluate to false  
if 后面与 != 都可以运行

- I/O

A scanf to read string will lead to buffer overrun problems.

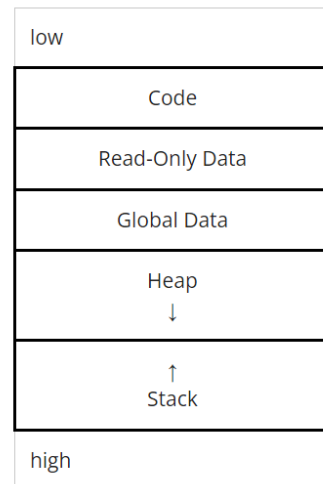
- String literal

def. "(string)" in an expression. (is initialization)  
read-only data section of memory

# 9. Dynamic Memory

## - Heap

def. final section in C memory model.



advantage of heap-allocated memory:

- 1) dynamic: the allocation size can be determined at run time
- 2) resizable: a heap allocation can be resized
- 3) duration: heap allocation persist until they are "freed"
- 4) safety: if memory runs out, it can be detected and handled properly.

## - dynamic memory

side effect: frees a. modifies "state" of the heap

most common use of dynamic memory: allocating space for arrays and structures

advantage: function can obtain memory that persists after obtained

side effects: modifies the "state" of heap.

## - malloc (memory allocation)

requires  $s$  bytes of memory from heap.

$\text{malloc}(s)$  → return pointer / NULL

## - realloc

resize array

- 步骤:
1. create new array.
  2. copy items
  3. free old

$\text{realloc}(p, \text{newsize})$   $O(n)$   
↑  
≥ 原 array in size

导致 dangling pointer 悬空指针

free a;

导致 memory leak

重复定义 a.

- merge sort  $(n \log n)$  ← worst case



## - uninitialized.

```
1 // This program demonstrates how memory returned
2 // from malloc is UNINITIALIZED
3
4 #include <stdlib.h>
5 #include <stdio.h>
6
7 struct posn {
8     int x;
9     int y;
10 };
11
12 int main(void) {
13     int *my_array = malloc(10 * sizeof(int));
14     my_array[3] = 2;
15     printf("the mystery value of my_array[0]: %d\n", my_array[3]);
16     printf("the mystery value of my_array[0]: %d\n", my_array[0]);
17     free(my_array);
18
19     struct posn *my_posn = malloc(sizeof(struct posn));
20     printf("the mystery values of my_posn: (%d, %d)\n", my_posn->x, my_posn->y);
21     free(my_posn);
22 }
```

分配 40 bytes memory 存到 \*my\_array

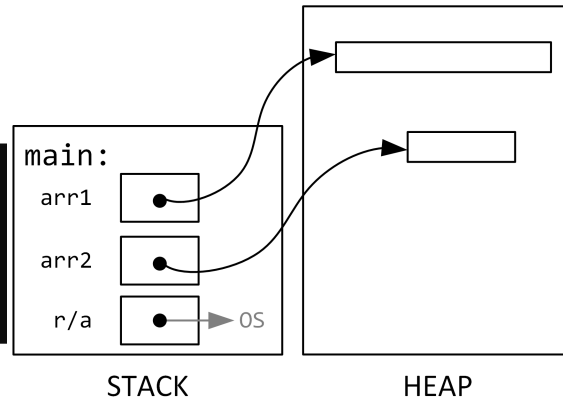
### Code Output

```
the mystery value of my_array[0]: 2
the mystery value of my_array[0]: -1094795586
the mystery values of my_posn: (-1094795586, -1094795586)
```

← initialized  
← uninitialized

## - Visualizing heap

```
int main(void) {
int *arr1 = malloc(10 * sizeof(int));
int *arr2 = malloc(5 * sizeof(int));
//...
}
```



out of memory

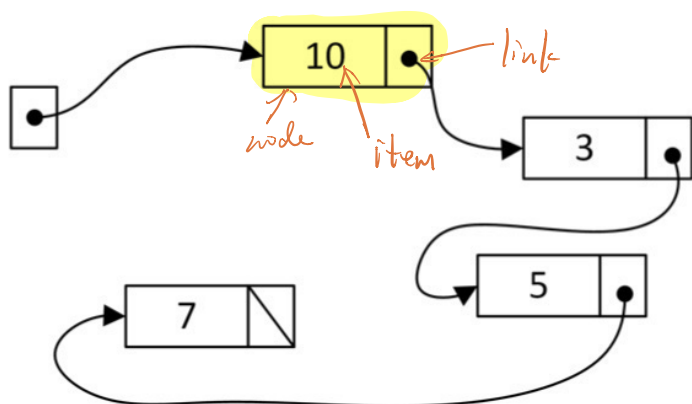
malloc ( sizeof(int) \* INT\_MAX )

free(p) return pointer p back to heap  
使 p 变为 dangling pointer (悬空指针)

## 10. ADT

- linked list

represented as a structure that contains a link (pointer) to the front node



recursive, data structure

```
1 struct llnode {
2     int item;
3     struct llnode *next;
4 };
5
6 struct llist {
7     struct llnode *front;
8 };
```

- Queue ADT

add-back

remove-front

front

is-empty

- Tree ADT

```
1 struct bstnode {
2     int item;
3     struct bstnode *left;
4     struct bstnode *right;
5     int count; // ****NEW
6 };
```

- dictionary ADT.

lookup

insert

remove

- ADT design.

combination of {  
primitives  
structures  
array  
linked lists  
trees  
graph

traditional collection ADT {  
stack  
queue  
sequence  
dictionary

}	dynamic array	sequenced data.	"specific position"
	sorted dynamic array	data are frequently searched.	X add remove
	linked list	sequenced. element 在头/尾 加/减	
	self-balancing binary search tree	unsequenced.	search / add remove

- generic ADT implementation.