

# On a Perplexing Polynomial Puzzle

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It seems rather surprising that any given polynomial  $p(x)$  with nonnegative integer coefficients can be determined by just the two values  $p(1)$  and  $p(a)$ , where  $a$  is any integer greater than  $p(1)$ . This result has become known as a “perplexing polynomial puzzle” in [2, 3]. Here, we address the natural question of what might be required to determine a polynomial with integer coefficients, if the condition that the coefficients be nonnegative is removed.

Let us analyze the original puzzle. Requiring that  $p(x) = c_0 + c_1x + \cdots + c_nx^n$  has nonnegative integer coefficients gives zero as a lower bound and  $p(1) = c_0 + c_1 + \cdots + c_n$  as an upper bound for the coefficients. Then if  $a > p(1)$ , writing  $p(a)$  as  $c_0 + c_1a + \cdots + c_na^n$  gives the unique base  $a$  representation of  $p(a)$ , so the nonnegative integer coefficients  $c_0, c_1, \dots, c_n$  of  $p(x)$  are completely determined by  $p(a)$ . The coefficients of  $p(x)$ , which serve as the base  $a$  digits, must fall in the appropriate set  $\{0, 1, \dots, a - 1\}$ , so we must choose  $a$  to be an upper bound of the coefficients.

If we allow negative coefficients and find a bound  $b$  on the coefficients so that  $-b \leq c_i \leq b$  for all  $i$ , then we wonder whether the coefficients are uniquely determined by the value of  $p(a) = c_0 + c_1a + \cdots + c_na^n$  for a suitable choice of  $a$ . Theorem 3 answers this question affirmatively, giv-

ing  $a = 2b + 1$  as a suitable choice. In effect, we ask whether  $p(a)$  has a unique base  $a$  representation if the *digits*  $c_i$  assume values from the set  $\{-b, -b + 1, \dots, 0, \dots, b - 1, b\}$ . Theorem 2 confirms the existence of such a representation.

## Nonstandard radix representations.

First we start with the uniqueness of the base  $a$  representation with a possibly nonstandard set of *digits*.

**Theorem 1** *Let  $a$  be a natural number and  $R = \{r_i \mid i = 0, 1, \dots, n\}$  be a set of integers such that  $r_i \not\equiv r_j \pmod{a}$  if  $i \neq j$ . If an integer  $z$  has a representation as a sum of powers of  $a$  with coefficients from  $R$ , then that representation is unique.*

*Proof:* Let  $z = \sum_{i=0}^m \lambda_i a^i = \sum_{j=0}^{m'} \mu_j a^j$  where all  $\lambda_i$  and all  $\mu_j$  are elements of  $R$ , and without loss of generality,  $m \leq m'$ . If  $\lambda_i = \mu_i$  for  $i = 0, 1, \dots, m$ , then clearly  $m = m'$  and we are done. Otherwise, assume  $s$  is the smallest index such that  $\lambda_s \neq \mu_s$ , and thus  $\lambda_i = \mu_i$  for  $i = 0, 1, \dots, s - 1$ . Consider the number  $u$  defined by

$$u = \frac{z - \sum_{i=0}^{s-1} \lambda_i a^i}{a^s} = \frac{z - \sum_{i=0}^{s-1} \mu_i a^i}{a^s} = \sum_{i=s}^m \lambda_i a^{i-s} = \sum_{j=s}^{m'} \mu_j a^{j-s}.$$

Now  $u \equiv \lambda_s \equiv \mu_s \pmod{a}$ . By assumption there exists at most one element in  $R$ , say  $r_s$ , such that  $r_s \equiv \lambda_s \pmod{a}$  and hence  $r_s = \mu_s = \lambda_s$ . So, there exists no smallest index  $s$  with  $\lambda_s \neq \mu_s$ , and the representation of  $z$  as a sum of powers of  $a$  with coefficients from  $R$  is unique. ■

Any integer  $z$  can be uniquely represented in base  $a$  as a sum of powers of  $a$  using coefficients from  $\{0, 1, \dots, a - 1\}$ . Next we show that unique representation as a sum of powers of  $a$  remains if we specify a different set of permissible coefficients, centered around 0.

**Theorem 2** *Let  $b$  be a natural number and let  $a = 2b + 1$ . Then every integer  $z$  can be uniquely written as  $z = \sum_{i=0}^m \lambda_i a^i$  where  $m \in \mathbb{N}$ ,  $\lambda_i \in \mathbb{Z}$ , and  $|\lambda_i| \leq b$  for each  $i = 0, \dots, m$ .*

*Proof:* If  $\sum_{i=0}^m \lambda_i a^i$  is the required representation of a nonnegative integer  $z$ , then  $\sum_{i=0}^m (-\lambda_i) a^i$  is the required representation of  $-z$ , so without loss of generality, we assume  $z$  is a nonnegative integer. Now  $z$  has a unique representation in base  $a$  as  $z = \sum_{i=0}^m \mu_i a^i$  where the integers  $\mu_i$  satisfy  $0 \leq \mu_i \leq 2b = a - 1$  for all  $i = 0, \dots, m$ . To shift the base  $a$  digits of  $z$  down by  $b$ , we add to  $z$  the base  $a$  number of equal length, all of whose digits are  $b$ , find the base  $a$  representation for this sum, then subtract the number to recover  $z$ . Specifically, consider  $z + \sum_{i=0}^m b a^i$ , which in base  $a$  is

$$z + \sum_{i=0}^m b a^i = \sum_{k=0}^{m'} \mu'_k a^k,$$

where  $0 \leq \mu'_k \leq 2b$  for each  $k$ . Observe that  $m' = m$  or  $m' = m + 1$  and  $\mu'_{m'} = 1$ . Hence

$$z = \sum_{k=0}^{m'} \mu'_k a^k - \sum_{i=0}^m b a^i = \begin{cases} \sum_{k=0}^{m'} (\mu'_k - b) a^k & \text{if } m' = m \\ a^{m'} + \sum_{k=0}^{m'-1} (\mu'_k - b) a^k & \text{if } m' = m + 1. \end{cases}$$

Letting  $\lambda_k$  be the coefficient of  $a^k$ , the leading coefficient  $\lambda_{m'}$  is either  $\mu'_{m'} - b$  or 1, and the trailing coefficients are  $\lambda_k = \mu'_k - b$ . Since  $0 \leq \mu'_k \leq 2b$ , we have  $|\mu'_k - b| \leq b$  and thus  $|\lambda_k| \leq b$  for all  $k = 0, 1, \dots, m'$ , as required. The uniqueness of the representation follows from Theorem 1. ■

We next see that, with a bound on the integer coefficients of a polynomial, the polynomial is uniquely determined by its value at one appropriately chosen input.

**Theorem 3** *Let  $p(x) = c_0 + c_1 x + \dots + c_n x^n = \sum_{i=0}^n c_i x^i$  be a polynomial with integer coefficients  $c_0, c_1, \dots, c_n$ . Assume  $b$  is a natural number such that  $|c_i| \leq b$  for all  $i = 0, 1, \dots, n$ . Let  $a = 2b + 1$ . Then the value  $p(a)$  uniquely determines the polynomial  $p(x)$ .*

*Proof:* Let  $p(x)$ ,  $a$ , and  $b$  be as in the statement of the theorem. By Theorem 2,  $p(a) = \sum_{i=0}^m c_i a^i$  can be uniquely written as  $\sum_{i=0}^m \lambda_i a^i$ , so  $p(a)$  uniquely determines the values of the coefficients  $c_i$  and thus of  $p(x)$ . ■

**Example:** Ask a friend to think of a polynomial  $p(x)$  with integer coefficients  $c_i$  satisfying  $|c_i| \leq 10$  for all  $i$ . We want to determine  $p(x)$  from just one value of  $p(x)$ . Applying Theorem 3, the bound on the coefficients is  $b = 10$ , so  $a = 2b + 1 = 21$  and we ask for the value  $p(21)$ . Suppose your friend's concealed polynomial was  $p(x) = 3x^4 - 5x^3 + 10x - 6$ . She would report  $p(21) = 537342$ .

Following the idea of the proof of Theorem 1 and using the Euclidean algorithm, we find the unique representation of  $p(21)$  as a linear combination  $\sum_{i=0}^m c_i 21^i$  of powers of 21 using integer coefficients  $c_i$  with  $|c_i| \leq 10$ . Since  $p(21) = 537342 = 25587(21) + 15 = 25588(21) - 6$ , we know  $c_0 = -6$  is the constant term of  $p(x)$ . We move to the next higher power of 21 and note that, besides the  $-6$  units,  $p(21)$  includes  $25588(21)$ , or 25588 in the 21's place. Now  $25588 = 1218(21) + 10$ , and since the remainder is between  $\pm 10$  inclusive, this remainder 10 must be  $c_1$ . Continuing in this manner,  $1218 = 58(21) + 0$ , and  $|0| \leq 10$ , so  $c_2 = 0$ ;  $58 = 2(21) + 16 = 3(21) - 5$ , and  $|-5| \leq 10$ , so  $c_3 = -5$ ; and finally,  $3 = 0(21) + 3$  so  $c_4 = 3$ . We have now recovered the polynomial  $p(x) = \sum_{i=0}^m c_i x^i = 3x^4 - 5x^3 + 0x^2 + 10x - 6$ .

## Refinements.

To apply Theorem 3, we need upper and lower bounds on the coefficients of the polynomial, which we then use to find a single bound  $b$  with  $|c_i| \leq b$  for each coefficient  $c_i$ . In the original puzzle, the requirement that the coefficients be nonnegative gave the lower bound on the coefficients as zero, while the upper bound could be subsequently deduced from  $p(1)$ . If we are only given a negative lower bound  $-b \in \mathbb{Z}$  for the integer coefficients of a polynomial  $p(x)$ , what else would be needed to deduce an upper bound for the coefficients? Giving one additional value of  $p(x)$  will no longer suffice: If  $p(a) = v$  is given

for  $a \neq 0$ , the possibilities for  $p$  include all polynomials

$$q(x) = v + (a^2 + a^4 + a^6 + \cdots + a^{2n}) - (x^2 + x^4 + x^6 + \cdots + x^{2n}),$$

and since the constant terms of these become arbitrarily large as  $n$  increases, no upper bound for the coefficients can be determined from only this information. If  $p(a) = v$  is given for  $a = 0$ , the possibilities for  $p$  include all polynomials  $q(x) = v + nx$ , and again the coefficients are unbounded. In the present case, assuming that a negative lower bound  $-b$  for the coefficients has been given, let us additionally assume that the leading coefficient is positive. Then, for  $a = 2b + 1$ ,  $a$  is positive, and, if  $n$  is the degree of  $p$ ,

$$p(a) \geq a^n + \sum_{i=0}^{n-1} (-b)a^i = a^n - b \sum_{i=0}^{n-1} a^i = a^n - b \frac{1 - a^n}{1 - a} = a^n + \frac{1}{2}(1 - a^n) = \frac{1}{2}(a^n + 1).$$

So, given  $p(a) = p(2b + 1)$ , the largest possible  $n$  satisfying  $(a^n + 1)/2 \leq p(a)$  is an upper bound on the degree of  $p(x)$ . This, in turn, can be used to find an upper bound for the coefficients  $c_i$  of  $p(x)$ . If  $p(x)$  has degree  $n$ , less than or equal to  $n_0$ , then  $p(a) = \sum_{i=0}^n c_i a^i$  and thus for  $0 \leq j \leq n$ , we have

$$c_j \leq c_j a^j = p(a) - \sum_{\substack{i=0 \\ i \neq j}}^n c_i a^i \leq p(a) + \sum_{\substack{i=0 \\ i \neq j}}^n b a^i \leq p(a) + \sum_{i=0}^{n_0} b a^i,$$

and the coefficients of  $p(x)$  have an upper bound. Now, as before, one more value of  $p(x)$ , chosen according to Theorem 3, will determine the whole polynomial. Thus, a lower bound on the integer coefficients, assuming the leading coefficient is positive, allows a polynomial to be determined by just two values at appropriate integers.

In [1] it was shown that a polynomial  $p(x)$  with nonnegative integer coefficients can actually be determined from sufficiently many digits of the value  $p(\pi)$ . Analogously, given a polynomial  $p(x)$  with integer coefficients bounded below by  $-b$  ( $b \in \mathbb{N}$ ) and with positive leading coefficient,  $p(x)$  can be determined from sufficiently many digits of  $p(t)$  for any transcendental number

$t > a = 2b + 1$ : If  $p(x)$  has degree  $n$ , we see that

$$\begin{aligned} p(t) &\geq t^n - b \sum_{k=0}^{n-1} t^k = t^n - b \frac{1 - t^n}{1 - t} \\ &\geq t^n - b \frac{1 - t^n}{1 - (2b + 1)} = \frac{1}{2}(t^n + 1). \end{aligned}$$

As before, the largest possible  $n$  with  $(t^n + 1)/2 \leq p(t)$  gives the maximal possible degree of  $p(x)$  and, as before, an upper bound on the coefficients of  $p(x)$ . Since the bound on the degree and the bounds on the integer coefficients of the polynomial in question narrow the possibilities to a finite number of polynomials, a sufficiently large finite number of digits of the value  $p(t)$  will determine the given polynomial.

It is easy to verify that an upper bound on the integer coefficients of a polynomial with negative leading coefficient can analogously be determined by either two of its values at appropriate integers or by one value at an appropriate transcendental number.

## References

- [1] F. Bornemann and S. Wagon, A perplexing polynomial puzzle revisited, *College Math. J.* **36** (2005) 288.
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